

```

/*+1=====*/
/*  MODULE                BOOL.C                */
/*=====*/
/*  FUNCTION      Boolean expression compiler  */
/*  */
/*  SYSTEM        Standard(ANSI / ISO) C.      */
/*               Tested on PC / MS DOS 5.0(MSC 600 A). */
/*  */
/*  SEE ALSO      Modules: GENERAL.H, STACK.H, ERROR.H/C, BOOL.H */
/*  */
/*  PROGRAMMER    Allan Dystrup.              */
/*  */
/*  COPYRIGHT(c)  Allan Dystrup, FEB.1991     */
/*  */
/*  VERSION       $Header: d:/cwk/kf/bool/RCS/bool.c 1.1 92/10/25 16:52:50 */
/*               Allan_Dystrup Exp Locker: Allan_Dystrup $ */
/*  -----*/
/*               $Log:   bool.c $ */
/*               Revision 1.1  92/10/25  16:52:50  Allan_Dystrup */
/*               Initial revision */
/*=====COMPILER STRUCTURE =====*/
/*
/* This module implements a simple compiler for evaluating boolean
/* expressions. The compiler is structured into two parts :
/* - a "front end" for transforming a boolean expression from infix form
/*   to postfix form (consisting of : SCANNER, PARSER, EMITTER);
/* - a "back end" for interpreting the postfix boolean expression using
/*   an abstract/software stack machine (consisting of : INTERPRETER).
/* The compiler "front end" and "back end" communicates using a common
/* symbol table (cf BOOL.H). The overall structure of the module is thus :
/*
/*   infix string---* SCANNER---* PARSER---* EMITTER---* postfix string
/*                   |                               * |
/*                   |                               | |
/*                   +-----* SYMBOL-----+ |
/*                               * | |
/*                               | | |
/*                   (FIND)-----+ |
/*                               * |
/*   evaluation result *-----INTERPRETER *-----+
/*
/*
/*
/*===== COMPILER FRONT END =====*/
/*
/* The syntax of the input string is defined by the following rules.
/*
/* 1. Boolean operator precedence, arity and associativity :
/*   operator(s)  precedence  arity and associativity
/*   -----
/*   NOT          1          unary
/*   AND          2          binary
/*   OR   XOR    3          binary, left associative
/*
/* 2. Context-free grammer (expressed in BNF:Bacus-Naur/Normal Form) :
/*   1. set of tokens(= terminal symbols) : { NOT, AND, OR, XOR, EOS, ID }
/*   NOT,AND,OR,XOR  The 4 basic Boolean operators
/*   EOS            End Of String
/*   ID            IDentifier
/*   2. set of nonterminals : { Z, E, T, F }
/*   Z             Distinguished start symbol
/*   E             Expression
/*   T             Term
/*   F             Factor

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/* 3. set of productions (context-free: one nonterminal on left side) */
/*      Z  ->  E EOS                                     */
/*      E  ->  E OR T | E XOR T | T                     */
/*      T  ->  T AND F | F                               */
/*      F  ->  ID | ( E ) | NOT F                       */
/*
/* We want to construct a recursive descent predictive parser, ie.
/* a top-down parser without backtracking. Furthermore we will require
/* the parser to work with 1 lookahead-character, ie. a LL(1) parser
/* (scanning input Left->right, using Leftmost derivations, 1 lookahead).
/* To make this possible, we have to transcribe the productions to get :
/* - no left recursion (thus eliminating infinite loops)
/* - no conflict between two right sides for any lookahead symbol
/* (ie. an unambiguous FIRST-set for any production) :
/*      Z  ->  E EOS                                     */
/*      E  ->  T E'                                       */
/*      E' ->  OR T E' | XOR T E' | epsilon
/*      T  ->  F T'                                       */
/*      T' ->  AND F T' | epsilon
/*      F  ->  ( E ) | NOT F | ID
/*
/* 3. Syntax-directed translation :
/* Combination of syntax analyzer and intermediate-code generator.
/* The chosen intermediate code is postfix (= "reverse polish"),
/* a notation in which the operators appear after their operands.
/* Postfix notation is a compact way of representing a parse tree,
/* with an unambiguous decoding dictated only by position and arity
/* (= number of operator arguments), - ie. parentheses are eliminated.
/* The interpretation of postfix is analogous to a bottom-up traversal
/* of the parse tree, - but may be performed by a simple stack machine.
/* We use a simple syntax directed definition (with semantic rules
/* defined by simple concatenation of translations) yielding a
/* corresponding simple translation scheme that can execute the
/* semantic actions (ie. emit intermediate code) during the parse.
/* Thus it is not nessecary to explicitly build the parse tree (syntax
/* analysys) before emitting the output (code ceneration); - hence
/* the concept "syntax-directed translation" !
/*
/*:-----:*/
/*:      Simple syntax directed definition      :      Translation scheme      :*/
/*:-----:*/
/*: Productions def.      : Semantic rules for      : Semantic actions in      :*/
/*: gramm. constructs : postfix=.p translation : procedural notation      :*/
/*: (infix)              : (infix to postfix)      : (generate postfix)      :*/
/*:-----:*/
/* Z  -> E EOS           : Z.p := E.p EOS           : Z -> E {terminate}      */
/* E  -> T E'            : E.p := T.p, E'.p         : E -> T E'               */
/* E' -> OR T E'         : E'.p := T.p, OR, E'      : E'-> OR T {print(OR)} E' */
/*      | XOR T E'       :      | T.p, XOR, E'      :      | XOR T {print(XOR)} E' */
/*      | epsilon        :      | epsilon           :      | epsilon           */
/* T  -> F T'            : T.p := F.p, T'.p         : T -> F T'               */
/* T' -> AND F T'        : T'.p := F.p, AND, T'p    : T'-> AND F {print(AND)} E' */
/*      | epsilon        :      | epsilon           :      | epsilon           */
/* F  -> ( E )           : F.p := E.p               : F -> ( E )              */
/*      | NOT F          :      | F.p, NOT          :      | NOT F {print(NOT)}   */
/*      | ID             :      | ID.p              :      | ID {print(ID.slot)}  */
/*:-----:*/
/*
/*
/*===== COMPILER BACK END =====*/
/*
/* The compiler back-end takes the front-end generated code (a postfix
/* string) as input, and evaluates it on the target machine (a simple
/* stack construction).
/*

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/* To be specific, the type of syntax-directed translation used by the */
/* compiler front end is based on an S-attributed definition with only */
/* synthesized attributes (more complex grammars require context info */
/* passed down the syntax tree in the form of inherited attributes). */
/*
/* The "internal form" emitted by the front end is an abstract syntax */
/* tree casted as a postfix string; More precicely, the syntax tree is */
/* a Directed Acyclic Graph : a "DAG", because the value index for an */
/* operator that occurs multiple times in a boolean expression, is */
/* unique (a reference to the symboltable) in the postfix string. */
/*
/* These characteristics allow a simple depth-first evaluation of the */
/* syntax tree (no dependency graph transformation of the tree needed). */
/* Thus the evaluation function may be implemented by a stack machine */
/* interpreting the postfix-string from left to right. */
/*
/*-1=====*/

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/*-----*/
/*                               INCLUDE HEADER-FILES                               */
/*-----*/

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/* Std.C header files */
#include <stdio.h>
#include <stddef.h>
#include <ctype.h>
#include <string.h>

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/* Project header files */
#include "general.h"
#include "error.h"
#include "stack.h"

```

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/* Module header file */
#define _BOOL_ALLOC
#include "bool.h"

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#define COPYRIGHT "Copyright (c) 1990 Allan Dystrup / KMD IS"
#define VERSION "v.1.0"

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/*-----*/
/*                               PRIVATE GLOBAL DATA                               */
/*-----*/

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/* Lexemes must by one char in range [0-127] */
/* Token      Lexeme      Semantics      */
#define OR     (BYTE) '/'   /* boolean OR   */
#define XOR    (BYTE) '%'   /* boolean XOR  */
#define AND    (BYTE) '&'   /* boolean AND  */
#define NOT    (BYTE) '^'   /* boolean NOT  */

#define LP     (BYTE) '('   /* left parenthesis */
#define RP     (BYTE) ')'   /* right parantesis */

#define QUOTE  (BYTE) ':'   /* quoted string   */
#define ESCAPE (BYTE) '\\ ' /* escape next char */

#define ID     (BYTE) 'I'   /* identifier      */
#define EOI    (BYTE) '\0'  /* End Of Input    */

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PRIVATE BYTE* yytext = (BYTE *) ""; /* Lexeme (not \0 terminated) */
PRIVATE int   yyleng = 0;          /* Lexeme length */
PRIVATE BYTE  aEmitStr[OUTMAX];    /* Output string from Emitter */

/*-----*/
/*                                     PRIVATE FUNCTION PROTOTYPES                                     */
/*-----*/

/* SCANNER */
PRIVATE BYTE
    bScan P((BYTE * pzStr));

/* PARSER */
PRIVATE void
    vAdvance P((BYTE token));
PRIVATE void
    vExpr    P((void));
PRIVATE void
    vTerm    P((void));
PRIVATE void
    vFactor  P((void));

/* EMITTER */
PRIVATE void
    vEmit    P((BYTE token));

/* SYMBOL */
PRIVATE int
    iSymInsert P((BYTE * pbStr, int len));

#ifdef MAIN
/*+2 MODULE BOOL.C =====*/
/*  NAME    00                main                               */
/*== SYNOPSIS =====*/
/* DESCRIPTION                Test driver for module bool.c
*-2*/
int
main(argc, apzArgv)
    WORD    argc;
    BYTE*   apzArgv[];
{
    int     i;
    BYTE    bPeek;
    FLAG    fResult;

    BYTE    aInfix[256]; /* input string : infix notation */
    BYTE    *pzPostfix; /* output string : postfix notation */

    /* 0: Echo input arguments */
    printf("\nargc=%d\n", argc);
    for (i = 0; apzArgv[i] != NULL; i++)
        printf("apzArgv[%d]=%s\n", i, apzArgv[i]);

    /* 1: Test SCANNER */
    printf("\nSCANNER ...\n");
    strcpy(aInfix, apzArgv[1]); /* save org. input string (aInfix) */

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printf("Infix : %s\n", aInfix);
bPeek = bScan(aInfix);
while (bPeek != EOI) {
    printf("Token : [%c] Lexeme : [yytext=%s yyleng=%d]\n", \
        bPeek, yytext, yyleng);
    bPeek = bScan("");
}

/* 2: Test PARSER */
printf("\nPARSER ...\n");
strcpy(aInfix, apzArgv[1]); /* restore org. input string (aInfix) */
pzPostfix = pzParse(aInfix);
printf("INFIX.. : %s\n", aInfix);
printf("POSTFIX : %s\n", pzPostfix);

/* 3: Test INTERPRETER */
printf("\nINTERPRETER ...\n");
for (i = 1; i < SYMMAX; i++) /* set up dymmy values for test */
    symtable[i].fValue = (i % 2 ? '\x01' : '\x00');
fResult = fInterpret(pzPostfix);
printf("Expression evaluates to : %s\n", (int) fResult ? "TRUE" : "FALSE");

return 0;

} /* END function main() */
#endif /* MAIN */

/*+2 MODULE BOOL.C =====*/
/** NAME 01 ***** SCANNER ******/
/*== SYNOPSIS =====*/
/*
 * DESCRIPTION Lexical analyzer
 * Linear analysis (= lexical analysis or scanning) of the input stream :
 * - The input string is read from left to right (using a "greedy algorithm")
 * - Characters are grouped into lexemes (sequences of characters with a
 * collective meaning, according to the defined grammer)
 * - For each lexeme matched by a defined pattern in the grammer the scanner
 * returns a unique type encoding (token), and assigns a pair of global
 * attribute values :
 * . "yytext" : a pointer to the lexeme (character string, not null term.)
 * . "yyleng" : an integer giving the length of the lexeme (in chars.)
 * The tokens returned by the scanner is used to direct the syntax analysis
 * in the parser-function.
 * - If the token is ID (ie. boolean vExpression identifier), there are many
 * possible lexemes for the same pattern, and the scanner then inserts the
 * lexeme in a global symboltable "symtable[]" for easy reference by the
 * later stages of the compilation process.
 * - The simple input string format does not warrant implementation of any
 * error recovery/transformation or any special buffering scheme.
*-2*/
BYTE
bScan(pzStr)
    BYTE* pzStr;
{
    static BYTE *index; /* Index = current pointer into input string */
    BYTE *p, *q, c;

    /* 1: Initialize & Skip leading whitespace */
    if (*pzStr)

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index = pzStr;

while (isspace(*index))
    ++index;

/* 2: Collect next token from pzStr; - break at valid token. */
/* (Guarantied loop termination at EndOfString EOS = \0) */
for (; TRUE; ++index) {
    yytext = index;          /* Point to first char in new token */
    yyleng = 1;

    switch (*index) {

        /* 2.1: One char operator */
        case EOI:             /* - terminate scan at EOI ! */
        case AND:
        case OR:
        case XOR:
        case NOT:
        case LP:
        case RP:
            return *index++;
            break;           /* Unreachable; - defensive programming */

        /* 2.2: Quoted string */
        case QUOTE:
            /* collect phrase lexeme */
            for (++index, ++yytext; *index != QUOTE; index++)
                if (*index == EOI)
                    vError(ELEX000, "Mangler slutmarkering af frase");

            yyleng = index - yytext;

            /* skip end-quote & return token */
            if (yyleng > 0) {
                index++;
                if (!iSymLookup(yytext, yyleng))
                    iSymInsert(yytext, yyleng);
                return ID;
            }
            break;          /* Empty phrase : ignore & get next token */

        /* 2.3: Identifier lexeme */
        default:
            /* Collect identifier lexeme, incl. escaped characters; */
            /* Break at whitespace or token, - unless escaped. */
            /* (Guarantied loop termination at EndOfString EOS = \0) */
            for (; TRUE; index++)
                /* Terminate at reserved character */
                if (isspace(c = *index) ||
                    c == AND || c == OR || c == XOR || c == NOT ||
                    c == LP || c == RP || c == QUOTE || c == EOI) {

                    /* unless previously ESCAPE'd (excluding EOI!) */
                    if (*(index - 1) == ESCAPE) {
                        for (p = index; p > yytext && *(p - 1) == ESCAPE; p--)
                            /* move p back over preceding ESCAPEs */;
                        if ((index - p) % 2 != 0 && c != EOI)
                            continue;
                    }
                    break;
                }
            }
            yyleng = index - yytext;
    }
}

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        /* Eliminate ESCAPE-character(s) from lexeme */
        for (p = yytext; p < index; p++)
            if (*p == ESCAPE) {
                for (q = p + 1; q < yytext + yyleng; q++)
                    *(q - 1) = *q;
                *(q - 1) = ' ';
                yyleng--;
            }

        /* Insert lexeme in symbol table */
        if (!iSymLookup(yytext, yyleng))
            iSymInsert(yytext, yyleng);

        return ID;
        break;          /* Unreachable; - defensive programming */

    } /* end switch */

} /* end for */

/* bScan() never returns implicitly by "falling out of scope" */
} /* END function bScan() */

```

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/*+2 MODULE BOOL.C =====*/
/** NAME    02      *****  PARSER ******/
/*== SYNOPSIS =====*/
/*
 * DESCRIPTION          Syntax analyzer
 * Hierachial analysis (= syntax analysis or parsing) where tokens are
 * grouped hierachially into grammatical phrases (a parse tree).
 * We construct a "predictive parser" (relying on our context-free grammer
 * with no left-recursion and unambiguous FIRST-set for all productions) :
 * For each lookahead symbol L, we :
 * - Determine the grammatical production G, given by "L in FIRST(G)";
 * - Advance the input stream to the next token (read next L);
 * - Call the procedure P implementing G, where P :
 *   . for each nonterminal calls a sub-procedure (with possible "recursive
 *     descent" calls to lower nonterminal procedures)
 *   . may call on a code generator to vEmit code according to the semantic
 *     actions indicated in a translation scheme for the grammer.
 *     The code is emitted into the global string "aEmitStr" (declared
 *     static for this module), the address of which is returned to the
 *     calling function as a "handle".
 * -2*/

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```
PRIVATE BYTE bLookahead;
```

```
PRIVATE void
vAdvance(bToken)
    BYTE bToken;
{
    if (bLookahead == bToken)          /* Match of expected token: */
        bLookahead = bScan((BYTE *) ""); /* advance to next token */
    else                                /* No match : */
        vError(ESYN000, (char *) yytext); /* fatal syntax error! */
}

```

```

/*----- pzParse -----*/
/* Z -> E {terminate} */
/*-----*/
BYTE*

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```

pzParse(pzStr)
  BYTE* pzStr;
{
  /* Check for valid argument */
  if (pzStr == NULL)
    vError(EARG000, "pzParse(NULL pointer)");

  /* Initialize SCANNER (scan first lookahead symbol) */
  bLookahead = bScan(pzStr);

  /* Start the ball rolling : top-down recursive descent of EXPRESSION */
  D(sprintf("parse[%c]\n", (char) bLookahead));
  vExpr();

  /* Terminate OK or ERROR */
  if (bLookahead == EOI) {
    vEmit(EOI);
    return aEmitStr;
  } else
    vError(ESYN000, (char *) yytext);
}

/*----- vExpr -----*/
/* E -> T E' */
/* E' -> OR T {print(OR)} E' */
/*   | XOR T {print(XOR)} E' */
/*   | epsilon */
/*-----*/
PRIVATE void
vExpr(void)
{
  BYTE      bOperator;

  D(sprintf(" vExpr[%c]\n", (char) bLookahead));
  vTerm(); /* T */

  while (TRUE) /* E' - perform recursively */
    switch (bLookahead) { /* (replaced by iteration) */

      case OR:
      case XOR:
        bOperator = bLookahead;
        vAdvance(bLookahead);
        D(sprintf(" vExpr[%c]\n", (char) bLookahead));
        vTerm();
        vEmit(bOperator);
        continue;

      default:
        /* E' finished */
        return;
    } /* END switch(bLookahead) */
}

/*----- vTerm -----*/
/* T -> F T' */
/* T' -> AND F {print(AND)} T' */
/*   | epsilon */
/*-----*/
PRIVATE void
vTerm(void)
{
  D(sprintf(" term[%c]\n", (char) bLookahead));
}

```



```

vFactor();          /* F */

while (TRUE)       /* T' - perform recursively */
  switch (bLookahead) { /* (replaced by iteration) */

    case AND:
      vAdvance(bLookahead);
      D(sprintf("  term[%c]\n", (char) bLookahead));
      vFactor();
      vEmit(AND);
      continue;

    default:
      /* F' finished */
      return;

  } /* END switch(bLookahead) */
}

/*----- vFactor -----*/
/* F -> ( E )          */
/*   | NOT E {print(NOT)} */
/*   | ID {print(ID.offset)} */
/*-----*/
PRIVATE void
vFactor(void)
{
  D(sprintf("  vFactor[%c]\n", (char) bLookahead));
  switch (bLookahead) {
    case LP:
      vAdvance(LP);
      vExpr();
      vAdvance(RP);
      break;

    case NOT:
      vAdvance(bLookahead);
      vFactor();
      vEmit(NOT);
      break;

    case ID:
      vEmit(ID);
      vAdvance(bLookahead);
      break;

    default:
      vError(ESYN000, (char *) yytext);

  } /* END switch(bLookahead) */
}

/*+2 MODULE BOOL.C =====*/
/** NAME    03      ***** EMITTER ******/
/*== SYNOPSIS =====*/
/*
 * DESCRIPTION          Intermediate code generator
 * Generates an intermediate representation of the input string (aInfix)
 * in the form of a postfix string ("reverse polish notation") suitable
 * for subsequent interpretation by an abstract (ie. software) stack machine.
 * The input string syntax (a boolean expression) is extremely simple

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* consisting of only identifiers and boolean operands; - No variables or
* control structures (apart from paranthetical nesting) are allowed, and
* hence the postfix syntax does not have to take into account assignments
* (l-values vs. r-values), control flow (jump instructions) et. al.
* The syntax of the generated postfix string is as follows :
* - for boolean operators we store their token value with the high bit on;
*   The token value for the boolean operators is identical to their lexeme
*   representation (range [0...127]), so the range for stored operators
*   is [128...255]. No operator attributes are stored in the symbol table.
* - for boolean identifiers we store their slot in the symbol table. Legal
*   values for symbol table slots are [1...SYMMAX], where SYMMAX <= 127
*   (to prevent collision with the value range for operators).
* - the postfix string is ended by a normal string terminator ('\0').
*   It is allocated as a global static string in this module, but it's
*   address is returned to the caller of parse as a handle.
*-2*/
PRIVATE void
vEmit(bToken)
    BYTE bToken;
{
    int        i;
    static BYTE *index = aEmitStr; /* Index = current pointer into output string */

    if ((index - aEmitStr) > OUTMAX)
        vError(ETAB000, "OUTMAX");

    switch (bToken) {

        case AND:          /* Boolean operator */
        case OR:
        case XOR:
        case NOT:
            D(sprintf("EMIT OP : [%c]\n", bToken | '\x80'));
            *index++ = (bToken | (BYTE)'\x80'); /* raise high bit */
            break;

        case ID:           /* boolean identifier */
            i = iSymLookup(yttext, yleng);
            D(sprintf("EMIT ID : [%c]-index[%d]->[%s]\n", (BYTE) i, i,
symtable[i].pzLexptr));
            *index++ = (BYTE) i;
            break;

        case EOI:
            D(sprintf("EMIT EOS: [NULL]\n"));
            *index++ = '\0';
            break;

        default:
            vError(ETOK000, ""); /* we should not end up here */
            break;
    }
}

/*+2 MODULE BOOL.C =====*/
/** NAME    04      ***** SYMBOL ******/
/*== SYNOPSIS =====*/
/*
* DESCRIPTION          Symbol-table manager
* Records input string identifiers in a symbol table, ie. a data structure
* containing a record for each identifier, w. fields for various attributes
* (in this case : a lexeme pointer and a boolean/flag value : "fValue").
* - The SCANNER enters the lexeme pointer into the symbol table for each

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*   new boolean vExpression identifier (the boolean attribute is undefined).
*   - The PARSER calls the EMITTER to lookup identifiers and enter their
*   symbol table slot into the intermediate code (ie. the postfix string).
*-2*/
static int  lastentry = 0;      /* last used position in symboltable */

PRIVATE BYTE lexemes[STRMAX]; /* allocate lexeme array */
static int  lastchar = -1;    /* last used position in lexemes */

PRIVATE int
iSymInsert(pbLex, len)
    BYTE* pbLex;
    int len;
{
/* Insert lexeme s as last entry in the symbol table
* Return slot for entry (or 0 if error).
*/

    if (lastentry + 2 >= SYMMAX)      /* room for array terminator */
        vError(ETAB001, "SYMMAX");
    if (lastchar + len + 2 >= STRMAX) /* room for string terminator */
        vError(ETAB002, "STRMAX");

    symtable[++lastentry].pzLexptr = &lexemes[++lastchar];
    strncpy((char *) symtable[lastentry].pzLexptr, (char *) pbLex, len);
    lastchar = lastchar + len;
    lexemes[lastchar++] = '\0';
    symtable[lastentry + 1].pzLexptr = NULL;

    D(sprintf("  INSERT :%tsymtable[%d] <- [%s]\n", \
        lastentry, symtable[lastentry].pzLexptr));
    return lastentry;
}

int
iSymLookup(pbLex, len)
    BYTE* pbLex;
    int len;
{
/* Determine whether there is an entry for lexeme s in the symbol table
* If yes, return slot in symbol table [1...SYMMAX], else return 0
*/
    int      p;

    for (p = lastentry; p > 0; p--)
        if (((int)strlen((char *) symtable[p].pzLexptr) == len) &&
            (strcmp((char *) symtable[p].pzLexptr, (char *) pbLex, len) == 0))
            return p;

    return 0;          /* no match */
}

void
vSymReset(void)
{
/* Reset boolean value for all lexemes in Symboltable to 'FALSE'
*/
    int      p;

    for (p = lastentry; p > 0; p--)
        symtable[p].fValue = FALSE;
}

```

```

/*+2 MODULE BOOL.C =====*/
/** NAME    05      ***** INTERPRETER ******/
/*== SYNOPSIS =====*/
/*
 * DESCRIPTION          Stack machine
 * Interprets a postfix representation of a boolean vExpression to T or F.
 *
 * The syntactic structure of the boolean vExpression is captured (by the
 * compiler frontend) as an "abstract syntax tree" in the condensed form of
 * a postfix string (cf. format specified by the EMITTER function). This
 * "internal form" is passed to the interpreter function as parameter.
 *
 * The actual evaluation is performed by an abstract stack machine, which
 * reads the postfix string from left to right (corresponding to a "depth
 * first evaluation" of the abstract syntax tree), and :
 * - pushes operand values onto the stack; An operand is represented in the
 *   postfix string as a "value number", ie. an index into the symbol table.
 *   The operand value (a boolean attribute "fValue") is retrieved from the
 *   symbol table and pushed on the stack
 * - performing operator actions :
 *   . popping boolean operands as required (cf. operator arity)
 *   . executing the boolean operation (using C-language boolean functions)
 *   . pushing the boolean fResult
 * until the whole postfix string is evaluated (ie. EOI is encountered), and
 * the result resides as the top stack-element.
 *
 * The abstract stack machine is implemented as a collection of general macro
 * definitions (in the header file "stack.h"). This implementation is chosen
 * for maximum execution speed, but the extensive use of macro "side effects"
 * (in the form of stack pointer arithmetic) and the use of unchecked stack-
 * operations does require some programming care.
*-2*/

```

```
stack_dcl(eval, int, SYMMAX);
```

```
FLAG
```

```
fInterpret(pzStr)
```

```
    BYTE* pzStr;
```

```
{
```

```
    BYTE          v;
```

```
    if (pzStr == NULL)
```

```
        vError(EARG001, "fInterpret(...NULL pointer...);
```

```
    stack_clear(eval);
```

```
    for (; *pzStr != EOI; pzStr++)
```

```
        if (*pzStr & '\x80')          /* Boolean operator ? */
```

```
            switch (*pzStr & '\x7F') { /* Mask off high bit */
```

```
                case AND:
```

```
                    D(printf("\tSTACK\t0->%c", stack_item(eval, 0) | 0x30));
```

```
                    D(printf("\t1->%c\n", stack_item(eval, 1) | 0x30));
```

```
                    v = (BYTE) (stack_item(eval, 0) & stack_item(eval, 1));
```

```
                    popn_(eval, 2);
```

```
                    push_(eval, v);
```

```
                    D(printf("\tAND"));
```

```
                    D(printf("\t0->%c\n", stack_item(eval, 0) | 0x30));
```

```
                    break;
```

```
                case OR:
```

```

        D(printf("\tSTACK\t0->%c", stack_item(eval, 0) | 0x30));
        D(printf("\t1->%c\n", stack_item(eval, 1) | 0x30));
        v = (BYTE) (stack_item(eval, 0) | stack_item(eval, 1));
        popn_(eval, 2);
        push_(eval, v);
        D(printf("\tOR"));
        D(printf("\t0->%c\n", stack_item(eval, 0) | 0x30));
        break;

    case XOR:
        D(printf("\tSTACK\t0->%c", stack_item(eval, 0) | 0x30));
        D(printf("\t1->%c\n", stack_item(eval, 1) | 0x30));
        v = (BYTE) (stack_item(eval, 0) ^ stack_item(eval, 1));
        popn_(eval, 2);
        push_(eval, v);
        D(printf("\tXOR"));
        D(printf("\t0->%c\n", stack_item(eval, 0) | 0x30));
        break;

    case NOT:
        D(printf("\tSTACK\t0->%c\n", stack_item(eval, 0) | 0x30));
        v = (BYTE) ((~pop_(eval)) & '\x01');
        push_(eval, v);
        D(printf("\tNOT"));
        D(printf("\t0->%c\n", stack_item(eval, 0) | 0x30));
        break;

    default:
        vError(ETOK001, "");    /* we should not end up here */
        break;
}
else {
    /* Boolean operand */
    D(printf("\tID\t%s\n", symtable[(int) *pzStr].pzLexptr));
    push_(eval, symtable[(int) *pzStr].fValue);
}

return ((pop_(eval)) == '\x00' ? FALSE : TRUE);
} /* END function fInterpret() */

/* END module BOOL.C */
/*=====*/
□

```