

```

/*=====*/
/* MODULE                                BM.C                                */
/*=====*/
/* FUNCTION    This module implements the Boyer-Moore algorithm
 *             for single substring search.
 *
 * SYSTEM      Standard C (ANSI/ISO).
 *             Tested on PC/MS DOS V.5.0.
 *
 * SEE ALSO    general.h  error.c/.h  bool.c/.h
 *
 * PROGRAMMER  Allan Dystrup
 *
 * COPYRIGHT   (c) Allan Dystrup
 *
 * VERSION     $Header: d:/cwk/kf/bm/RCS/bm.c 1.1 92/10/25 17:02:46
 *             Allan_Dystrup Exp Locker: Allan_Dystrup $
 *             -----
 *             $Log:    bm.c $
 *             Revision 1.1 92/10/25 17:02:46 Allan_Dystrup
 *             Initial revision
 *
 * REFERENCES
 *
 * [1] Donald E. Knuth, James H. Morris, Jr. and Vaughan R. Pratt :
 *     "Fast pattern matching in strings",
 *     SIAM J. COMPUT. Vol. 6, No. 2, June 1977.
 *
 * [2] Robert S. Boyer andj. Strother Moore :
 *     "A fast string searching algorithm",
 *     Communications of the ACM, October 1977, Vol. 20, No. 10.
 *
 * [3] Donald M. Sunday :
 *     "A very fast substring search algorithm",
 *     Communications of the ACM, August 1990, Vol. 33, No. 8.
 *
 * USAGE
 *

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*===== SUBSTRING SEARCH =====
*
* An elementary problem in information retrieval is searching for a specific
* substring or "pattern" (length p) in a larger string of text (length t),
* where pattern and text are composed of characters from a fixed alphabet
* (length a).
*
*
*                               SF, Straight Forward
* The SF (or "brute force") method consists of first aligning the start of
* the pattern with the start of the text, and then performing a left-right
* scan of the pattern and a char. by char. comparison with the text.
* In case of a mismatch the pattern is shifted right one step, and the text-
* pointer is "backed up" to the new position of the pattern start.
* The time complexity of SF is :
*   setup time : zero
*   worst-case : quadratic in  $O(p*t)$ 
*                 (ex. find "aaaab" in "aaaaa...aaaaab",
*                 ie. match whole pattern FOR EACH char. in the text)
*   average    : linear in  $O(t)$ , with const. factor  $[1..p]$  near 1.
*                 (mismatch usually expected at first char. comparison,
*                 in which case there is NO backtracking in the text)
*
*
*                               KMP, Knuth-Morris-Pratt - ref. [1]
* The KMP method performs a left-right pattern scan like SF, but eliminates
* "backing up" the text-pointer at character mismatch by first setting up a
* transition graph including each character in the pattern string.
* The graph is represented by a "next[]" array with an index for each pattern
* character giving the new position IN THE PATTERN STRING for backing up THE
* PATTERN POINTER in case of character mismatch, thereby shifting the pattern
* forward by "delta" characters to align the already MATCHED textstring
* (=pattern suffix) with the first leftmost matching pattern prefix/infix.
* The KMP delta shift precalculation thus utilizes any reoccurring substrings
* in the pattern string, and also eliminates backtracking the textpointer
* (if mismatch and no delta shift : start new scan from next textchar)
* KMP time complexity :
*   setup time : linear in  $O(p)$ 
*   worst-case : linear in  $O(t)$ , with const. factor = 1
*                 (ex. find "abcde" in "fffff...fffabcde")
*   average    : linear in  $O(t)$ , with const. factor  $[1/p..1]$  near 1.
*                 (mismatch usually expected at first char. comparison,
*                 in which case KMP delta shift IS NOT utilized at all).
*
*
*                               BM, Boyer-Moore - ref. [2]
* The original BM method performs a "REVERSE" pattern scan (ie. right-left,
* unlike SF and KMP), using two predefined shift functions :
* - delta1 : For each char 'c' in the alphabet, find the rightmost position
*             (i) of c in the pattern (i=0 if c not in pat).  $\Delta_1(c) = p-i$ .
* - delta2 : Analogous to the KMP delta function for reoccurring substrings
*             in the pattern : for each pattern suffix find first leftward
*             reoccurrence of it (infix or prefix) in the pattern.
* In case of mismatch BM uses the larger value of delta1 and delta2 to shift
* the pattern forward :
* - either  $\Delta_1(c)$  : to the rightmost matching char. 'c' in the pattern,
*   however if  $\Delta_1(c)$  is to the right of the mismatch-position, then
*   the pattern is shifted right by 1,
* - or  $\Delta_2(c)$  : to the first leftward pattern prefix/infix matching the
*   already matched textstring, cf. KMP algorithm.
* and restarting the scan from the pattern end,
* BM time complexity :
*   setup time : linear in  $O(a+p)$ 
*   worst-case : - using delta1 alone... : quadratic in  $O(p*t)$ 
*                 - using delta1 & delta2 : linear in  $O(t)$ 
*   average    : sub-linear in  $O(t/p)$ , w. const. factor  $[1/p..1]$  near  $1/p$ .
*                 (mismatch usually expected at first char. comparison,
*                 in which case BM delta shifts ARE fully utilized).
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*===== CHOICE OF ALGORITHM =====
*
* The only advantage of SF is that it doesn't require any setup time. For
* very simple search situations (small and heterogenous pattern and text)
* SF may be worth considering, but as a general basis for bibliographic
* search SF is totally inadequate.
*
* KMP has an advantage when "backing up" the textpointer is inconvenient, or
* when searching for a highly self-repetitive pattern in a self-repetitive
* text. Under these circumstances the linear worst-case search time makes it
* more attractive than SF, but in most actual applications (as for instance
* bibliographic search for a single word) KMP is not likely to be signifi-
* cantly faster than SF.
*
* BM is our candidate for a bibliographic one-word search algorithm. The
* "sublinear" search time of delta1 makes BM more than three times as fast
* as SF or KMP in most practical situations, and with the addition of delta2
* we get an algorithm with the same desirable worst-case performance as KMP.
* With a small modification of the original BM algorithm it is even possible
* to get a simpler and faster implementation that does not depend on a
* specific pattern scan order and thus may be coded to avoid backtracking!
* The advantage of BM increases with the length of the search pattern (p) and
* of the alphabet (a). A bibliographic search for a long keyword or phrase
* in a text over the full (255 char) alphabet is an ideal application for BM.
*
* BM however has one drawback : it can not easily be extended to search for
* multiple patterns in a text. For this type of application it is better to
* use the technique applied by KMP : building a transition graph (also
* called a "Finite State Automaton") to perform the search. This is precisely
* the scope of the Aho-Corasic algorithm for multiple string search (see
* module ac.c).
*
*===== AN IMPROVED BM IMPLEMENTATION =====
*
* An improved implementation of the BM algorithm, - cf. ref. [3]
*
* Delta1.
* First note that the pattern string always shifts right by at least one
* char. Hence the char. in the text string IMMEDIATELY PAST THE END of the
* pattern string must be involved for testing at the next pattern position.
* Thus delta1 can be computed for the whole alphabet to be the index of the
* first leftward occurrence of each char. FROM THE END of the pattern string.
* This slightly modified delta1 has the following big advantage :
* the shift defined by the new delta1(c) is an ABSOLUTE SHIFT relating to
* the text character immediately after the pattern (always >= 1), hence
* (1) delta1 MAY BE USED "STAND ALONE" to code a "quick and clean" BM
* search (as opposed to using the original BM delta1 SHIFT, which was
* relative to the mismatch position, requiring a test and shift by
* max(1, delta2) when the pos. of mismatch was to the left of delta1.
* (2) delta1 does NOT DEPEND ON A SPECIFIC PATTERN SCAN ORDER (as opposed
* to the original BM, which required a strict right-left "backtracking"
* textpointer)
*
* Delta2
* For any specific order of scanning the pattern string, one can define a
* delta2 shift similar to the KMP delta (left-right scan) or the BM delta2
* (right-left scan), thus avoiding the undesirable "quadratic worst case
* behaviour" of a BM delta1 stand-alone solution (but at the expense of a
* substantially larger setup time).
*
* Delta1 vs. Delta2.
* Empirical evidence indicates a "break even" between delta1 standalone and
* delta1+delta2 at a pattern length of about 10-15 characters (in biblio-
* graphic search), - ie. a search for a short keyword is best performed by
* a delta1 standalone algorithm, while searching for a longer word or phrase
* often justifies the initial setup time for delta2.
*
*/

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/*=====*/
/*                               Includes                               */
/*=====*/
/* Standard ANSI/ISO header files */
#include <stdio.h>
#include <limits.h>
#include <string.h>
#include <stdlib.h>

/* Module header file */
#define _BM_ALLOC                /* bm.c (Boyer-Moore) header file */
#include "bm.h"

/*=====*/
/*                               Typedefs                               */
/*=====*/
typedef struct patScanElem { /* Struct. of element in scan ordered pattern */
    int      loc;           /* Location of the character c in the pattern */
    BYTE     c;            /* Value of character in pattern at location */
} sPAT;

/*=====*/
/*                               Global data objects                     */
/*=====*/
PRIVATE BYTE *pzPat;        /* Ptr. to pattern string */
PRIVATE int iPatLen;        /* Length of pattern string */
PRIVATE int *piDelta1;     /* Ptr. to Delta1 shift array for all chars */
PRIVATE float *pfScanOrd;  /* Ptr. to scan priority array for chars */
PRIVATE sPAT *psOrdPat;    /* Ptr. to scan ordered elements of pattern */
PRIVATE int *piDelta2;     /* Ptr. to Delta2 shift array for pat. chars */
PRIVATE enum scanType eScanOrder = eScanUD; /* Initially undefined */

/*=====*/
/*                               Function Prototypes                     */
/*=====*/
PRIVATE int iShCompMS(sPAT * psPat1, sPAT * psPat2);
PRIVATE int iShCompOM(sPAT * psPat1, sPAT * psPat2);
PRIVATE int iShFind(int iPatLoc, int iLShift);

/*=====*/
/*                               Trace macros                             */
/*=====*/
/* The option of using an arbitrary scan order in the improved BM algorithm
 * requires a rather convoluted procedure for setting up the piDelta2-shifts.
 * I have defined the following macros to allow tracing and verification of
 * the BM table-setup. The notation of the macro parameters uses the follo-
 * wing "shorthand":
 *   _P   : pointer to start of zero-terminated pattern string
 *   _p   : pointer for scanning the pattern string
 *   _s   : pointer to array of scan order priority for characters
 *   _OP  : pointer to start of list of ordered structures 'patScanElem'
 *   _o   : pointer for scanning the _OP list
 *   _arr : array
 *   _len : length of array or string
 *   _i   : scratch index variable
 */
int      _i;

#define DUMPD1(_arr, _len)                                     \
D(      printf("\n\nDUMP of BM table "#_arr"[] ... \n");      \

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    for (_i=0; _i < _len; _i++) \
        printf(" %x[%c]:%04.1f%c", \
            _i, _i, (float) _arr[_i], !(_i%6) ? '\n' : ' '); );

#define DUMPMS(_typ,_o,_p,_s,_len) \
D(    printf("\n\nDUMP of "#_typ" scan ordered string "#_o" ...\n"); \
    for(_i=0; _i < _len; _i++) \
        printf("\t"#_o "[%02d]="#_p "[%02d]=%c  "#_s "[%02d]=%04.1f\n", \
            _i, _o[_i].loc, _o[_i].c, _i, _s[_o[_i].loc] ); );

#define DUMPOM(_typ,_o,_p,_s,_len) \
D(    printf("\n\nDUMP of "#_typ" scan ordered string "#_o" ...\n"); \
    for(_i=0; _i < _len; _i++) \
        printf("\t"#_o "[%02d]="#_p "[%02d]=%c  "#_s "[%c]=%04.1f\n", \
            _i, _o[_i].loc, _o[_i].c, _o[_i].c, _s[_o[_i].c] ); );

#define DUMPD2(_typ,_arr,_len) \
D(    printf("\n\nDUMP of "#_typ" BM table "#_arr"[] ...\n"); \
    for(_i=0; _i < _len; _i++) \
        printf("\t"#_arr "[%02d]=%02d\n", _i, _arr[_i]); );

#define DUMPSH(_typ,_len,_OP,_o,_P,_j) \
D(    printf("\t  iLShift=%d OP[%02d]=P[%02d]=%c P[%02d]=%c\t"#_typ"\n", \
        _len, _o-_OP, _o->loc, _o->c, _j, (_j>=0 ? _P[_j] : '*') ); );

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#ifdef MAIN
/*=====*/
/*                               Main (Module Testdriver)                               */
/*=====*/
#define MAX_LINE      256

int
main(int iArgc, BYTE ** ppzArgv)
/* Input is a Boolean search expression "str" provided on the command line.
 * The Boolean expression is parsed by function pzParse (in bool.c) to :
 * - a suffix representation of the Boolean logic tree (pzPostfix).
 * - a set of keywords K = {k1, k2, ... ,kn} to search for (symtable[]).
 * In module bm.c we perform a simple search for a keyword or phrase using
 * the improved Boyer-Moore algorithm. The use of pzParse to scan/parse the
 * single keyword or phrase is thus somewhat "overkill" (we don't actually
 * use pzPostfix in this context), but the driver may easily be extended to
 * run other more advanced search algorithms as for instance the Aho-Corasic
 * search for multiple substrings in a text.
 */
{
    BYTE      *pbArg;          /* Char in argument string */
    int       eScanOrder = eScanUD; /* Type of pattern scan order */
    FILE      *fileInpFd;     /* File descriptor for input file */
    BYTE      pzInpBuffer[MAX_LINE]; /* String buffer for input file */
    BYTE      *pzPostfix;     /* Ptr. to postfix boolean logic */
    BYTE      *pbNL;         /* Prt. to newline character */
    FLAG      fMatch;        /* Flag for match Pattern/Text */
    int       i;             /* Scratch index variable */

    /* 1: Parse command line option */
    while (--iArgc && (**++ppzArgv)[0] == '-')
        for (pbArg = ppzArgv[0] + 1; *pbArg != '\0'; pbArg++)
            switch (toupper(*pbArg)) {
                case 'M':      /* Force a Maximal Shift search */
                    eScanOrder = eScanMS;
                    break;
                case 'O':      /* Force an Optimal Mismatch search */
                    eScanOrder = eScanOM;
                    break;
                case 'Q':      /* Force a Quick Search */
                    eScanOrder = eScanQS;
                    break;
                default:       /* Switch must be M, O, or Q */
                    printf("\n\n--- allowed switches : [MOQ] ! ---\n\n");
                    exit(EXIT_FAILURE);
            }
    }

    /* 2: Parse Bool search expression to postfix string and keyword list */
    /* -- (Use only first entry in keyword list : symtable[1] for search) */
    D(printf("\n\n===== BUILDING NEW BM =====\n\n"));
    D(puts("PARSE INPUT STRING ..."));
    D(printf("\t%s\n", *ppzArgv));
    pzPostfix = pzParse(*ppzArgv++);
    iArgc--;

    for (i = 1; symtable[i].pzLexptr; i++)
        /* count symbol table entries */;
    if (--i > 1 || !symtable[i].pzLexptr) {
        printf("\n\n--- exactly one search phrase, please! ---\n\n");
        exit(EXIT_FAILURE);
    }

    /* 3: Initialize BM-search */
    if (eScanOrder == eScanUD)
        eScanOrder = (strlen(symtable[1].pzLexptr) <= 15 ? eScanQS : eScanMS);
}

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vBuildBM(eScanOrder, symtable);

/* 3: Run BM-search */
fileInpFd = (iArgc ? fopen(*ppzArgv, "r") : stdin);
while (fgets(pzInpBuffer, MAX_LINE, fileInpFd)) {

    if (pbNL = strchr(pzInpBuffer, '\n'))/* Replace newline w. string term. */
        *pbNL = '\0';

    if (fMatch = fRunBM(pzInpBuffer)) {
        D(sprintf("\nOUTPUT ... \n>\t"));
        puts(pzInpBuffer);
    }
}

/* 3: Terminate BM search */
/* Return status to the parent process: 1 if match, 0 if none */
vDelBM();
return (fMatch ? 1 : 0);
}
#endif /* #ifdef MAIN */

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/*=====*/
/*          vBuildBM          */
/*=====*/
void
vBuildBM(enum scanType type, struct entry * syntab)
{
/* Construct the BM delta shift table(s) from pattern string pzPat.
 *
 *          1 Delta1 table : piDelta1[]
 * The delta1 table-setup requires two steps :
 *
 * 1.1 First initialize the table to (iPatLen+1) for all characters in the
 *     alphabet.
 *
 * 1.2 Then reset the table entries for each char in the pattern to the
 *     rightmost position in the pattern of that char.
 *
 *          2 Delta2 table : piDelta2[]
 * The delta2 shift table piDelta2[] is built in two steps :
 *
 * 2.1 First build the ordered pattern psOrdPat from any ordering spe-
 *     cification; The chosen scan-order (set up in array pfScanOrd) may be :
 *     - forward, giving a piDelta2[] = KMP delta
 *     - reverse, giving a piDelta2[] = original BM delta2
 *     - maximal shift (MS), using the max. left shift to sort psOrdPat
 *     - optimal mismatch (OM), using char-frequency to sort psOrdPat
 *     - any other order suitable for the actual implementation
 *     For bibliographic search we choose to implement the MS scan order.
 *
 * 2.2 Then construct piDelta2[] in two steps, using the following procedure:
 *     If mismatch at scan order position i, ie psOrdPat[i].c !=
 *     text[ScanStart+psOrdPat[i].loc], then shift ('overlay') the pattern
 *     left to position 'iLShift', so
 *     2.2.1 psOrdPat[0].c ... psOrdPat[i-1].c ALL MATCH aligned characters
 *           in the shifted pat; - iLShift is the max. value in the range
 *           0...psOrdPat[0].loc (= the minimal leftshift to match i chars)
 *     2.2.2 psOrdPat[i].c DOES NOT match the aligned character
 *           pzPat[ psOrdPat[i].loc - iLShift ]; - this may require
 *           repeated left shifting of the pattern.
 *     Then delta2 table piDelta2[i] = (psOrdPat[0].loc - iLShift),
 *     ie. the required right shift to align reoccurring characters in the
 *     pattern (cf. 2.2.1) excluding the mismatched character (cf. 2.2.2)
 *     with already matched characters in the text.
 */
BYTE      *p;          /* Ptr. to scan through pattern string */
sPAT      *o;          /* Ptr. to scan through ordered pattern */
int        iLShift;    /* Left shift for matching reoccurring chars */
int        i, j;       /* Scratch index variables */

/* Define global data objects of module */
pzPat = syntab[1].pzLexptr; /* Pattern, -scanned by module bool.c */
iPatLen = strlen((char *) pzPat); /* Pattern length */
eScanOrder = type; /* Requested type of pattern scan order */
pfScanOrd = NULL; /* Defined in comp. functions below */

/* ----- Build delta1 table ----- */

/* Allocate the Delta1 shift table piDelta1[] */
if (!(piDelta1 = (int *) malloc(sizeof(int) * (UCHAR_MAX + 1))))
    vError(EMEM005, "vBuildBM");

/* 1.1 Initialize the piDelta1[] table for all characters in the alphabet */
for (i = 0; i <= UCHAR_MAX; i++) /* mismatch -> 1. char after pattern */
    piDelta1[i] = iPatLen + 1;

/* 1.2 Reset piDelta1[] entries for each of the pattern characters */
for (p = pzPat; *p; p++) /* match -> last position in pattern */

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        piDelta1[*p] = iPatLen - (p - pzPat);
DUMPD1(piDelta1, UCHAR_MAX);

/* ----- Build delta2 table ----- */

if (eScanOrder != eScanQS) {

    /* Alloc. structures for psOrdPat and Delta2 shift table piDelta2[] */
    if (!(psOrdPat = (sPAT *) malloc(sizeof(sPAT) * (iPatLen + 1)))
        || !(piDelta2 = (int *) malloc(sizeof(int) * (iPatLen + 1))))
        vError(EMEM006, "vBuildBM");

    /* 2.1 */
    /* Build the scan-ordered pattern from the pfScanOrd array */
    /* using qsort to sort the pattern chars after max left shift */
    for (i = 0, p = pzPat, o = psOrdPat; i <= iPatLen; ++i, ++p, ++o) {
        o->loc = i;
        o->c = *p;          /* \0 for i == iPatLen */
    }

    switch (eScanOrder) {
        case eScanMS:      /* Maximal Shift search */
            qsort(psOrdPat, iPatLen, sizeof(sPAT), iShCompMS);
            DUMPMS(MS - sorted, psOrdPat, pzPat, pfScanOrd, iPatLen);
            break;
        case eScanOM:      /* Optimal Msimatch search */
            qsort(psOrdPat, iPatLen, sizeof(sPAT), iShCompOM);
            DUMPOM(OM - sorted, psOrdPat, pzPat, pfScanOrd, iPatLen);
            break;
        /* case tNEW : insert call to new scan order sorting here & */
        /* ----- place code of the scan order comp.func below */
        default:
            /* Must be a defined type of search ! */
            vError(EARG002, "vBuildBM");
    }

    /* 2.2.1 */
    /* First init piDelta2[] with the minimum matching left shift */
    /* so psOrdPat[0].c ... psOrdPat[i-1].c match aligned characters */
    for (piDelta2[0] = iLShift = i = 1; i < iPatLen; ++i) {
        iLShift = iShFind(i, iLShift);
        piDelta2[i] = iLShift;
    }
    DUMPD2(initialized, piDelta2, iPatLen);

    /* 2.2.2 */
    /* Then for each piDelta2[i] check that : psOrdPat[i].c != */
    /* pzPat[psOrdPat[i].loc-iLShift]; - if not, repeat search for a */
    /* matching left shift until this condition. */
    for (i = 0; i < iPatLen; ++i) {
        iLShift = piDelta2[i];          /* get initial matching shift */
        o = psOrdPat + i;
        D(sprintf("\n piDelta2[%02d], INITIAL LSHIFT : %d\n", i, iLShift));

        while (iLShift < iPatLen) {

            /* Require left shift INSIDE pattern! */
            if ((j = (psOrdPat[i].loc - iLShift)) < 0) {
                DUMPSH(Ignore, iLShift, psOrdPat, o, pzPat, j);
                break;
            }

            /* Require current char NOT MATCH after shift! */
            if (psOrdPat[i].c != pzPat[j]) {

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        DUMPSH(OK, iLShift, psOrdPat, o, pzPat, j);
        break;
    }

    /* If match, scan iteratively for next matching shift */
    DUMPSH(RESET, iLShift, psOrdPat, o, pzPat, j);
    ++iLShift;
    iLShift = iShFind(i, iLShift);
}
piDelta2[i] = iLShift;          /* set final shift */
}
DUMPD2(corrected, piDelta2, iPatLen);
}
/* end if eScanOrder != eScanQS */
}

/*-----*/
/*                               iShCompMS                               */
/*-----*/
int
iShCompMS(sPAT * psPat1, sPAT * psPat2)
{
/* Compare 2 elements of the scan ordered pattern according to a precalcula-
 * ted table 'pfScanOrd' of each element-character's maximal left shift.
 */
    int        dsh;          /* Delta shift value */
    int        i, j;        /* Scratch index variables */

/* First call: Build the pfScanOrd array for each char in pattern pzPat */
/* of the minimal left shift to the last position of the char in pzPat */
/* (or to the 1. character of pzPat if no match). */
    if (pfScanOrd == NULL) {
        if (!(pfScanOrd = (float *) malloc(sizeof(float) * (iPatLen + 1))))
            vError(EMEM007, "iShCompMS");
        for (i = 0; i < iPatLen; ++i) {
            for (j = i - 1; j >= 0 && pzPat[j] != pzPat[i]; --j);
            pfScanOrd[i] = (float) (i - j);
        }
        DUMPMS(unsorted, psOrdPat, pzPat, pfScanOrd, iPatLen);
    }

/* ascending order */
    dsh = (int) (pfScanOrd[psPat2->loc] - pfScanOrd[psPat1->loc]);
    return (dsh ? dsh : psPat2->loc - psPat1->loc);
}

/*-----*/
/*                               iShCompOM                               */
/*-----*/
int
iShCompOM(sPAT * psPat1, sPAT * psPat2)
{
/* Compare 2 elements of the scan ordered pattern according to a precalculated
 * table 'pfScanOrd' of each element-character's frequency in the alphabet. */

    int        i;
    float      fq;

/* First call: Build the pfScanOrd array for each char in the alphabet */
/* of freq. of the chars in the targetlanguage based on the alphabet. */
    if (pfScanOrd == NULL) {
        if (!(pfScanOrd = (float *) malloc(sizeof(float) * (UCHAR_MAX + 1))))
            vError(EMEM008, "iShCompOM");
        for (i = 0; i <= UCHAR_MAX; ++i)

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        pfScanOrd[i] = 0.1f;

/* English text alphabet frequency */
pfScanOrd['e'] = 11.1f;
pfScanOrd['a'] = 8.9f;
pfScanOrd['i'] = 7.8f;
pfScanOrd['r'] = 7.4f;
pfScanOrd['t'] = 7.1f;
pfScanOrd['o'] = 6.9f;
pfScanOrd['n'] = 6.8f;
pfScanOrd['s'] = 5.6f;
pfScanOrd['l'] = 5.5f;
pfScanOrd['c'] = 4.5f;
pfScanOrd['u'] = 3.6f;
pfScanOrd['m'] = 3.2f;
pfScanOrd['d'] = 3.2f;
pfScanOrd['p'] = 3.1f;
pfScanOrd['h'] = 2.9f;
pfScanOrd['g'] = 2.4f;
pfScanOrd['b'] = 2.3f;
pfScanOrd['y'] = 2.0f;
pfScanOrd['f'] = 1.5f;
pfScanOrd['w'] = 1.1f;
pfScanOrd['k'] = 1.1f;
pfScanOrd['v'] = 1.0f;
pfScanOrd['x'] = 0.3f;
pfScanOrd['j'] = 0.2f;
pfScanOrd['z'] = 0.2f;
pfScanOrd['q'] = 0.2f;

/* Danish text alphabet frequency */
/* - to be defined - */

DUMPD1(pfScanOrd, UCHAR_MAX);
DUMPOM(unsorted, psOrdPat, pzPat, pfScanOrd, iPatLen);
}

/* descending order */
fq = pfScanOrd[psPat1->c] - pfScanOrd[psPat2->c];
return (int) (fq ? (fq > 0 ? 1 : -1) : psPat2->loc - psPat1->loc);
}

/*-----*/
/*                               iShFind                               */
/*-----*/
int
iShFind(int i, int iLShift)
{
/* Return value of the next valid leftward shift 'iLShift' of the first 'i'
 * chars in psOrdPat, such that all chars from psOrdPat[i-1].c to
 * psOrdPat[0].c in the scan ordered pattern matches the chars in the pattern
 * pzPat after a left shift of 'iLShift' (ie. we match reoccurring chars in the
 * pattern at a distance of iLShift positions).
 */
    sPAT    *o;          /* Ptr. to scan ordered pattern */
    int      j;          /* Scratch index variable */

    D(sprintf("\n\t MATCHSHIFT psOrdPat[%02d...00] :\n", i - 1));
    for (; iLShift < iPatLen; ++iLShift) {

        for (o = psOrdPat + i - 1; o >= psOrdPat; o--) {
            /* in range 0...psOrdPat[j].loc */
            if ((j = (o->loc - iLShift)) < 0) {
                DUMPSH(Ignore, iLShift, psOrdPat, o, pzPat, j);
                continue;
            }
        }
    }
}

```

```
    if (o->c != pzPat[j]) {          /* mismatch at position j */
        DUMPSH(Mismatch, iLShift, psOrdPat, o, pzPat, j);
        break;
    }

    DUMPSH(MATCH, iLShift, psOrdPat, o, pzPat, j); /* matched one more */
}

if (o < psOrdPat)          /* matched all psOrdPat[0]...psOrdPat[i] */
    break;
}

D(printf("\t LSHIFT : %d\n", iLShift));
return iLShift;
}
```

```

/*=====*/
/*                                     vRunBM                                     */
/*=====*/
FLAG
fRunBM(BYTE * T)
{
/*           psOrdPat<----o---->psOrdPat+iPatLen
 * Scan ordered pattern : [xxxxxxxxxxx]
 *
 *           pzPat<----p--->pzPat+iPatLen
 * Pattern :             [xxxxxxxxxxx]
 *
 * Text      : [yyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyyy]
 *           T          S<----t--->S+iPatLen          T+Tlen
 */
  BYTE *p;          /* Ptr. to scan through pattern */
  sPAT *o;         /* Ptr. to scan ordered pattern */
  BYTE *S = T;     /* Ptr to shift the pattern through text */
  BYTE *t;        /* Ptr. to scan through text */
  int Tlen = strlen((char *) T); /* Length of text */
  int iD1;        /* Deltal table shift value. */
  int iD2;        /* Delta2 table shift value. */

  switch (eScanOrder) {
    case eScanQS:
      /* ----- Perform left-right Quick Search ----- */

      D(puts("\nTRACE of BM Quick Search ...\n"));

      /* while enough text is still left ... */
      while (S + iPatLen <= T + Tlen) {

        D(printf("Pat:\t%s\n", pzPat));
        D(printf("Txt:\t%s\n", S));

        /* scan through pattern string while match in text */
        D(printf("\tMATCH:["));
        for (p = pzPat, t = S; *p && *p == *t; p++, t++)
          D(printf("%c", *t));

        /* test for end-of-string */
        if (*p) { /* no : MISMATCH */
          D(printf("\n\tFAIL.: [%c<>%c]", *p, *t));
          iD1 = piDeltal[* (S + iPatLen)]; /* get deltal */
          D(printf("\n\tSHIFT:piDeltal[%c]=%03d\n", *(S + iPatLen),
iD1));

          S += iD1; /* shift pattern */
        }
        else { /* yes : MATCH */
          D(puts("\n\tCOMPLETE!\n"));
          return (TRUE); /* might return position (1+S-T) */
        }
      }
      return (FALSE); /* no match in total text, return FALSE(0) */
      break; /* unreachable,- defensive programming */

    case eScanMS:
    case eScanOM:
      /* ----- Perform arbitrary scan order search ----- */

      D(puts("\nTRACE of BM ScanOrder Search ...\n"));

      /* while enough text is still left ... */
      while (S + iPatLen <= T + Tlen) {

        D(printf("Pat:\t%s\n", pzPat));

```

```

D(printf("Txt:\t%s\n", S));

/* scan through ordered pattern while match in text */
D(printf("\tMATCH:["));
for (o = psOrdPat; o->c && o->c == *(S + o->loc); o++)
    D(printf("%c", *(S + o->loc)));

/* test for end-of-string */
if (o->c != 0) { /* no : MISMATCH */
    D(printf("\n\tFAIL.:pos.%03d [%c<>%c]\n",
            o->loc, o->c, *(S + o->loc)));
    iD1 = piDelta1[*(S + iPatLen)]; /* get delta1 */
    iD2 = piDelta2[o - psOrdPat]; /* get delta2 */
    D(printf("\tSHIFT:piDelta1[%c] =%03d\n", *(S + iPatLen),
            iD1));

    D(printf("\t      piDelta2[%03d]=%03d\n", o - psOrdPat, iD2));
    D(printf("\t      MAX      =%03d\n", (iD1 > iD2 ? iD1 : iD2)));
    S += (iD1 > iD2 ? iD1 : iD2); /* use max for pattern shift
*/

}
else { /* yes : MATCH */
    D(puts("\n\tCOMPLETE!\n"));
    return (TRUE); /* might return position (1+S-T) */
}
}
return (FALSE); /* no match in total text, return FALSE(0) */
break; /* unreachable,- defensive programming */

default: /* Must be a defined type of search ! */
    vError(EARG003, "vRunBM");
}
}

```

```

/*=====*/
/*                                     vDelBM                                     */
/*=====*/
void
vDelBM(void)
{
    switch (eScanOrder) {
        case eScanMS:
        case eScanOM:
            free(piDelta2);
            free(psOrdPat);
            free(pfScanOrd);
            /* !fall through! */
        case eScanQS:
            free(piDelta1);
            break;
        default:
            /* Must be a defined type of search ! */
            vError(EARG004, "vDelBM");
    }

    D(puts("\nData structures deallocated ..."));
}

/* END of module BM.C                                     */
/*=====*/

```